

Amendments to the Specification:

Please replace the paragraph beginning at page 4, line 15 with the following amended paragraph:

FIG. 1 shows the system architecture for performing the scheduling method of the present invention, wherein the number of flows in the system is assumed to be N. The packet from flow i is processed by the present scheduling method and output through the output queue 11. The output queue 11 has a plurality of windows 111, each window 111 having a size of W. The flows 1, 2, ...N have weights of $W_1^*, W_2^*, \dots, W_N^*$, respectively. Based on the size W of each window 111, the weight W_i^* of each flow i can be normalized to be $w_i = W \times W_i^* / (W_1^* + W_2^* + \dots + W_N^*)$.

Please replace the paragraph beginning at page 5, line 8 with the following amended paragraph:

If step S11 determines that the credit c_i of flow i is larger than the size of packet P_i , the packet P_i is placed into one of the windows 111, for example the kth window (window k of FIG. 3) 111, pointed by the window index d_i for being output through the output queue 11 (step S13). In step S14, the size of the packet P_i is subtracted from the credit c_i , and the window index w_i and the credit c_i are updated; i.e., the updated window index w_i and the credit c_i are written into the table 12 (step S14).

Please replace the paragraph beginning at page 5, line 15 with the following amended paragraph:

The packets placed in the windows 111 of the output queue 11 are sequentially output. As shown by the packet departure process in FIG. 1, when all packets have been pushed out and the window 111 is empty, the table 12 is updated to have its initial values.